# Queens and Pawns on Square Boards 

## Handout for the 2015 MOVES Talk

"The Maximum Queens Problem on a Rectangular Board"

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August 4, 2015

Each diagram below shows the minimum known number of pawns whose placement on the given board will allow the placement of the indicated number of mutually nonattacking queens.

## $5 \times 5$ boards




8 queens and 7 pawns


9 queens and 16 pawns

## $6 \times 6$ boards



6 queens and 0 pawns


8 queens and 3 pawns


7 queens and 1 pawn


9 queens and 6 pawns

## $7 \times 7$ boards



7 queens and 0 pawns


8 queens and 1 pawn


9 queens and 2 pawns


11 queens and 6 pawns


13 queens and 10 pawns


10 queens and 4 pawns


12 queens and 8 pawns


14 queens and 12 pawns


15 queens and 20 pawns


16 queens and 33 pawns

## $8 \times 8$ boards



10 queens and 2 pawns


11 queens and 3 pawns


14 queens and 9 pawns


13 queens and 7 pawns


15 queens and 11 pawns


## Sources

## The $n$-queens problem

The patterns with no pawns are well-known solutions to the $n$ queens problem. For more information on this problem, see Bell, J. \& Stevens, B. (2009). A survey of known results and research areas for $n$-queens. Discrete Math. 309, no. 1, 1-31.

## The $n+k$-queens problem

The patterns where the number of queens equals the order of the board plus the number of pawns ( 7 queens with 1 pawn on a $6 \times 6$ board; 8 queens with 1 pawn and 9 queens with 2 pawns on a $7 \times 7$ board; and 9 queens with 1 pawn, 10 queens with 2 pawns, and 11 queens with 3 pawns on an $8 \times 8$ board) are examples of solutions to the " $n+k$-queens problem". See http://npluskqueens.info for more information on this problem.

## Binary integer programming

To find each of the other patterns in this handout, I formulated the problem of finding the minimum number of pawns needed to allow placement of $Q$ mutually nonattacking queens on a given size board as a binary integer programming problem and submitted the problem to the NEOS server (http://www.neos-server.org/neos/). Here is the formulation of the "minimum pawns for $Q$ queens problem" on an $n \times n$ board:

Suppose we have an $n \times n$ board with rows labeled $0, \ldots, n-1$ and columns labeled $0, \ldots, n-1$.

For $i=0, \ldots, n-1$ and $j=0, \ldots, n-1$, let $q_{i j}=1$ if and only if there is a queen in row $i$ and column $j$ and $q_{i j}=0$ otherwise. Also, for each $i$ and $j$, let $p_{i j}=1$ if and only if there is a pawn in row $i$ and column $j$ and $p_{i j}=0$ otherwise.

Minimize $\Sigma_{i, j} p_{i j}$ (the number of pawns on the board) subject to the following constraints:

1. $\sum_{i, j} q_{i, j}=Q$. (There are exactly $Q$ queens on the board.)
2. For $0 \leq i \leq n-1$ and $0 \leq j \leq n-1,0 \leq q_{i j}+p_{i j} \leq 1$. (A queen cannot occupy the same space as a pawn.)
3. For $0 \leq i \leq n-1$ and $0 \leq j_{1}<j_{2} \leq n-1, \sum_{j=j_{1}}^{j_{2}}\left(q_{i j}-p_{i j}\right) \leq 1$. (No segment of a row has two queens without a pawn between them.)
4. For $0 \leq j \leq n-1$ and $0 \leq i_{1}<i_{2} \leq n-1, \sum_{i=i_{1}}^{i_{2}}\left(q_{i j}-p_{i j}\right) \leq 1$. (No segment of a column has two queens without a pawn between them.)
5. a) For $0<s \leq n-1$ and $0 \leq j_{1}<j_{2} \leq s, \sum_{j=j_{1}}^{j_{2}} q_{s-j, j}-p_{s-j, j} \leq 1$.
b) For $n \leq s<2 n-2$ and $s-(n-1) \leq j_{1}<j_{2} \leq n-1, \sum_{j=j_{1}}^{j_{2}} q_{s-j, j}-p_{s-j, j} \leq 1$. (No segment of a "sum diagonal" - the squares $(i, j)$ for which $i+j=s$ for some constant $s$ - has two queens without a pawn between them.)
6. a) For $n-1>d \geq 0$ and $0 \leq j_{1}<j_{2} \leq(n-1)-d, \sum_{j=j_{1}}^{j_{2}} q_{d+j, j}-p_{d+j, j} \leq 1$
b) For $-1 \geq d>-(n-1)$ and $-d \leq j_{1}<j_{2} \leq n-1, \sum_{j=j_{1}}^{j_{2}} q_{d+j, j}-p_{d+j, j} \leq 1$ (No segment of a "difference diagonal" - the squares $(i, j)$ for which $i-j=d$ for some constant $d$ - has two queens without a pawn between them.)

With a few changes, we get the integer programming formulation of the "minimum pawns for maximum queens problem" on an $m \times n$ board:

Suppose we have an $m \times n$ board with rows labeled $0, \ldots, m-1$ and columns labeled $0, \ldots, n-1$, and $m \leq n$. ( $m \leq n$ is an arbitrary choice which makes items 5 and 6 easier to state.)

For $i=0, \ldots, m-1$ and $j=0, \ldots, n-1$, let $q_{i j}=1$ if and only if there is a queen in row $i$ and column $j$ and $q_{i j}=0$ otherwise. Also, for each $i$ and $j$, let $p_{i j}=1$ if and only if there is a pawn in row $i$ and column $j$ and $p_{i j}=0$ otherwise.

Minimize $\Sigma_{i, j} p_{i j}$ subject to the following constraints:

1. $\sum_{i, j} q_{i, j}=\left\lceil\frac{m}{2}\right\rceil\left\lceil\frac{n}{2}\right\rceil$. (There are exactly $\left\lceil\frac{m}{2}\right\rceil\left\lceil\frac{n}{2}\right\rceil$ queens on the board.)
2. For $0 \leq i \leq m-1$ and $0 \leq j \leq n-1,0 \leq q_{i j}+p_{i j} \leq 1$. (A queen cannot occupy the same space as a pawn.)
3. For $0 \leq i \leq m-1$ and $0 \leq j_{1}<j_{2} \leq n-1, \sum_{j=j_{1}}^{j_{2}}\left(q_{i j}-p_{i j}\right) \leq 1$. (No segment of a row has two queens without a pawn between them.)
4. For $0 \leq j \leq n-1$ and $0 \leq i_{1}<i_{2} \leq m-1, \sum_{i=i_{1}}^{i_{2}}\left(q_{i j}-p_{i j}\right) \leq 1$. (No segment of a column has two queens without a pawn between them.)
5. a) For $0<s \leq m-1$ and $0 \leq j_{1}<j_{2} \leq s, \sum_{j=j_{1}}^{j_{2}} q_{s-j, j}-p_{s-j, j} \leq 1$.
b) For $m \leq s<m+n-2$ and $s-(m-1) \leq j_{1}<j_{2} \leq n-1, \sum_{j=j_{1}}^{j_{2}} q_{s-j, j}-p_{s-j, j} \leq$ 1.
(No segment of a "sum diagonal" - the squares $(i, j)$ for which $i+j=s$ for some constant $s$ - has two queens without a pawn between them.)
6. a) For $m-1>d \geq 0$ and $0 \leq j_{1}<j_{2} \leq(n-1)-d, \sum_{j=j_{1}}^{j_{2}} q_{d+j, j}-p_{d+j, j} \leq 1$
b) For $-1 \geq d>-(m-1)$ and $-d \leq j_{1}<j_{2} \leq n-1, \sum_{j=j_{1}}^{j_{2}} q_{d+j, j}-p_{d+j, j} \leq 1$
(No segment of a "difference diagonal" - the squares $(i, j)$ for which $i-j=d$ for some constant $d$ - has two queens without a pawn between them.)
